

The (expressive) power of graph learning Floris Geerts (University of Antwerp)

## Course

* Is about recent advances in graph learning.
* With an emphasis on the expressive power of learning methods.
* Self-contained (too some extent).
* Mostly high-level, but also low-level, so basically all levels.
* Not all methods or related works are covered.
* Will not report experiments...


## About the speaker

* Background in mathematics, database :. theory* and expressive power of query languages.
* Since 2018, expressive power of linear algebra.
* Natural move to the study of expressive power of graph neural networks.


## Outline

* Graph learning and expressive power
* Message Passing Neural Networks
- Boosting power:
* Feature augmentation
* Subgraphs
* Higher-order message-passing
Ask Questions



## Why learning on graphs?

## Graphs are everywhere!




## Graphs: One definition to rule them all

* Graph $G=\left(V_{G}, E_{G}, L_{G}\right)$ with
* Vertex set $V_{G}$
* Edge set $E_{G} \subseteq V_{G}^{2}:=V_{G} \times V_{G}$
* Vertex labels: $L_{G}: V_{G} \rightarrow \Sigma$



## Graphs: One definition to rule them all

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*Vertex labels: $L_{G}: V_{G} \rightarrow \Sigma$
Vertex features $\mathbb{R}^{d}$



## Adjacency matrix representation

* Graph $G=\left(V_{G}, E_{G}, L_{G}\right)$ can also be represented by adjacency matrix $A_{G}$ and feature matrix $F_{G}$
* Let $n=\left|V_{G}\right|$ be the number of vertices. Let $v, w \in[n]:=\{1, \ldots, n\}$.
adjacency matrix $\quad A_{G} \in \mathbb{R}^{n \times n}:(v, w) \mapsto \begin{cases}1 & (v, w) \in E_{G} \\ 0 & \text { otherwise }\end{cases}$
feature matrix $\quad F_{G} \in \mathbb{R}^{n \times d}: v \mapsto L_{G}(v)$
* Assumes an ordering on the vertices.


## Graph learning


$\mathbb{R}^{d}$


Classical ML
$\mathscr{G}=$ all graphs
$\mathbb{Y}=$ output space

## Graph learning



$$
\mathscr{G}=\text { all graphs }
$$

$$
\mathbb{Y}=\text { output space }
$$

## Embeddings

$$
\begin{aligned}
\mathscr{G} & =\text { all graphs } \\
\mathscr{V} & =\text { all vertices } \\
\mathbb{Y} & =\text { output space }
\end{aligned}
$$

* Graph embedding: $\xi: \mathscr{G} \rightarrow \mathbb{Y}$
* Vertex embedding: $\xi: \mathscr{G} \rightarrow(\mathscr{V} \rightarrow \mathbb{Y})$
* p-Vertex embedding: $\xi: \mathscr{G} \rightarrow\left(\mathscr{V}^{p} \rightarrow \mathbb{Y}\right)$


## Graph embeddings

* Graph embedding: $\xi: \mathscr{G} \rightarrow \mathbb{V}$
* Graph classification/regression



## Vertex embeddings

* Vertex embedding: $\xi: \mathscr{G} \rightarrow(\mathscr{V} \rightarrow \mathbb{Y})$
* Vertex classification/regression. For example, prediction of subject of papers.



## p-Vertex embeddings

*p-Vertex embedding: $\xi: \mathscr{G} \rightarrow\left(\mathscr{V}^{p} \rightarrow \mathbb{Y}\right)$

* For example, 2-vertex embeddings: link prediction



## Graph learning tasks



## Applications

* Vertex classification: categorise online user/items, location amino acids (protein folding, alpha fold)
* Link prediction: knowledge graph completion, recommender systems, drug side effect discovery
* Graph classification: molecule property, drug discovery
* Subgraph tasks: traffic prediction



## Applications

* Vertex classifica categorise online user/items, location aming (protein folding, alpha fold) GRAPH
* Link, LEARNING HAS ph completion, recomm BECOME KEY o side effect discovery DATA
* Graph clas COMPONENT ecule property, drug discovery
* Subgrap sks: traffic r liction



## Graph learning

* We want to learn an unknown embedding $\Xi: \mathscr{G} \rightarrow\left(\mathscr{V}^{p} \rightarrow \mathbb{V}\right)$


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What does this mean???

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What does this mean???

* The embedding $\Xi$ is partially revealed by means of a training set

$$
\mathscr{T}:=\left\{\left(G_{1}, \mathbf{v}_{1}, y_{1}\right), \ldots,\left(G_{\ell}, \mathbf{v}_{\ell}, y_{\ell}\right)\right\} \subseteq \mathscr{G} \times \mathscr{V}^{p} \times \mathbb{V}
$$

## Graph learning

* We want to learn an unknown embedding $\Xi: \mathscr{G} \rightarrow\left(\mathscr{V}^{p} \rightarrow \mathbb{Y}\right)$


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\boldsymbol{\uparrow} \\
\Xi\left(G_{1}, \mathbf{v}_{1}\right) \quad \Xi\left(G_{\ell}, \mathbf{v}_{\ell}\right)
\end{gathered}
$$

## Training sets



Graph classification

(cora, paper, topic)
Vertex classification

(social, $p_{x}, p_{y}$, yes $/$ no)
Link prediction

## Graph learning: hypothesis class

* We want to find the best model consistent with training set $\mathscr{T}$


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What does this mean???


## Graph learning: hypothesis class

* We want to find the best model consistent with training set $\mathscr{T}$ $\uparrow$ What does this mean???
- Models are selected from an hypothesis class $\mathscr{H}$
* In the graph setting $\mathscr{H}$ consists of embeddings


## Hypothesis classes



## Hypothesis classes

\section*{MPNN <br> GSN 2-IGN

\section*{PPGN

## PPGN $\times 1000$

Graphormer GATs CayleyNet
CWN $\delta-k-$ GNNs GIN GCNs
ChebNet
k-IGNs
GraphSage
Dropout GNN
k-GNNs
$\mathscr{H}$

## Explosion



## Graph learning

Classical embedding methods depend on representation E.g., think of MLP on vector representation of flattened adjacency matrix
Invariant
Embedding method

$$
\mathscr{G}=\text { all graphs }
$$

$$
\mathbb{Y}=\text { output space }
$$

## A desired property: Invariance

* Embeddings should be invariant, that is, independent of the chosen graph representation.
* Invariance is defined in terms of graph isomorphisms.

* The mapping $\pi$ is a bijective vertex function satisfying $\left(v, v^{\prime}\right) \in E_{G} \Longleftrightarrow(\pi(v), \pi(w)) \in E_{H}$ also $\underline{L_{G}(v)=L_{H}(\pi(v)) \text { must hold. }}$


## Invariant embeddings

for all $\pi, G$ and $\mathbf{v} \in V_{G}^{p}: \xi(G, \mathbf{v})=\xi(\pi(G), \pi(\mathbf{v}))$

Isomorphism

$(1,4)$ and $(\mathrm{B}, \mathrm{C})$ have same embedding in $\mathbb{Y}$
We typically assume invariant embedding methods (unless said otherwise)

## Graph learning: ERM

Best one! $\xi$

* Given training set $\mathscr{T}$ and hypothesis class $\mathscr{H}$
* Empirical risk minimisation:

Find embedding $\xi$ in $\mathscr{H}$ which minimises empirical loss

$$
\frac{\left.\frac{1}{\ell} \sum_{i=1}^{\ell} \operatorname{loss}\left(\xi\left(G_{i}, \mathbf{v}_{i}\right), y_{i}\right)\right)}{\mathbf{q}}
$$

Loss function is a mapping from $\mathbb{Y} \times \mathbb{Y} \rightarrow \mathbb{R}$

## Loss functions

* Choice depends on learning task (regression, classification,...)
* L1: $\operatorname{loss}\left(y_{\text {predicted, }}, y_{\text {true }}\right):=\left|y_{\text {predicted }}-y_{\text {true }}\right|$
* L2: loss $\left(y_{\text {predicted }}, y_{\text {true }}\right):=\left(y_{\text {predicted }}-y_{\text {true }}\right)^{2}$
- (Binary) cross entropy: $\operatorname{loss}\left(y_{\text {predicted }}, y_{\text {true }}\right):=y_{\text {true }} \log \left(y_{\text {predicted }}+\left(1-y_{\text {true }}\right) \log \left(1-y_{\text {predicted }}\right)\right.$


## Graph learning

* Graph learning systems solve ERM using back propagation and gradient descent...

$$
\left.\hat{\xi}: \arg \min _{\xi \in \mathscr{H}} \frac{1}{\ell} \sum_{i=1}^{\ell} \operatorname{loss}\left(\xi\left(G_{i}, \mathbf{v}_{i}\right), y_{i}\right)\right)
$$

## Graph learning

* Graph learning systems solve ERM using back propagation and gradient descent...


Our focus will be on the expressive power of hypothesis classes

## Expressive power

* Which embeddings can be expressed by embeddings in $\mathscr{H}$ ?
* Which embeddings can be approximated by embeddings in $\mathscr{H}$ ?
* Which inputs can be separated/distinguished by embeddings in $\mathscr{H}$ ?


## Notions of expressivity I

* Let $\Xi: \mathscr{G} \rightarrow\left(\mathscr{V}^{p} \rightarrow \mathbb{Y}\right)$ be a $p$-vertex embedding and let $\mathscr{C}$ be a subset of $\mathscr{G}$
$\mathscr{H}$ can $\mathscr{C}$-express $\Xi$ if there exists a $\xi \in \mathscr{H}$ such that for all $G \in \mathscr{C}, \mathbf{v} \in V_{G}^{p}$ :
$\xi(G, \mathbf{v})=\Xi(G, \mathbf{v})$
$\mathscr{H}$ can $\mathscr{C}$-approximate $\Xi$ if for any $\epsilon>0$
there exists a $\xi_{\epsilon} \in \mathscr{H}$ such that for all $G \in \mathscr{C}, \mathbf{v} \in V_{G}^{p}:\left\|\xi_{\epsilon}(G, \mathbf{v})-\Xi(G, \mathbf{v})\right\| \leq \epsilon$


## Notions of expressivity II

Separation/distinguishing power of $\mathscr{H}$

$$
\rho(\mathscr{H}):=\{(G, \mathbf{v}, H, \mathbf{w}) \mid \forall \xi \in \mathscr{H}: \xi(G, \mathbf{v})=\xi(H, \mathbf{w})\}
$$

* All pairs of inputs that cannot be separated by any embedding in $\mathscr{H}$


## Distinguishing power

- Strongest power: $\mathscr{H}$ powerful enough to detect non-isomorphic graphs
* Weakest power: $\mathscr{H}$ cannot differentiate any two graphs



## Distinguishing power

* Allows for comparing different classes of embeddings methods!

$$
\rho(\text { methods } 1) \subseteq \rho(\text { methods } 2)
$$

Methods1 is more powerful than Methods2
Methods 2 is bounded by Methods 1 in power

$$
\rho(\text { methods } 1)=\rho(\text { methods } 2)
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Both methods are as powerful

* Allows for comparing embedding methods with algorithms, logic, ...


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## Expressive power in ML community

* Focus has been on distinguishing power of classes $\mathscr{H}$ of embedding methods.
* Goal is to characterise $\rho(\mathscr{H})$ in a way to sheds light on what graph properties a learning method can detect/use.
* We see an example shortly for $\mathscr{H}=$ the class of Message-Passing Neural Networks (MPNNs)


## Expressive power in ML community

* Search for increase in expressive power has led to surge of new methods of graph learning.
* Despite theoretical underpinning... still a bit of alchemy to find the right method...


Li hico! la panarrece l'age murr la sumease


We will gradually fill in this landscape with recent graph learning methods

## Questions?



The most popular type of $G \mathcal{N N s}$

## A little history

Sll

## A little history



## Message passing neural networks

A class of invariant vertex and graph embedding methods


## Idea behind MPNNs: Neighbour aggregation



Every vertex defines a computation graph

Neural networks


## MPNNs: Vertex embedding

$\xi(G, v):=\xi^{(L)} \circ \xi^{(L-1)} \circ \cdots \circ \xi^{(0)}(G, v)$
Message Passing Layers

$$
\begin{aligned}
& \xi^{(0)}(G, v):=\text { Hot-one encoding of label of vertex } v \\
& \xi^{(t)}(G, v):=\operatorname{Upd}^{(t)}\left(\xi^{(t-1)}(G, v), \operatorname{Agg}^{(t)}\left(\left\{\left\{\xi^{(t-1)}(G, v), \xi^{(t)}(G, u) \mid u \in N_{G}(v)\right\}\right\}\right)\right)
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$$



## MPNNs: Vertex embedding



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Message Passing between $v$ and its neighbours $u \in N_{G}(v)$


## MPNNs: Vertex embedding



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& \begin{array}{c}
\text { Message Passing between } v \text { and its } \\
\text { Ueighbours } u \in N_{G}(v)
\end{array} \\
& \text { Update and aggregate function contain }
\end{aligned}
$$

## MPNNs: Graph embedding

$$
\frac{\rho(G):=\rho \circ \xi^{(L)} \circ \xi^{(L-1)} \circ \ldots \circ \xi^{(0)}(G, v)}{\text { Readout }}
$$

$$
\begin{aligned}
& \rho(G):=\operatorname{Readout}\left(\left\{\left\{\xi^{(L)}(G, v) \mid v \in V_{G}\right\}\right\}\right) \\
& \text { Has learnable parameters }
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$$

Typical choices for update, aggregate and readout: Multilayer Perceptrons

## MPNNs: Graph embedding

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& \rho(G):=\operatorname{Readout}\left(\left\{\left\{\xi^{(L)}(G, v) \mid v \in V_{G}\right\}\right\}\right) \\
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& \text { Aggregation over all vertices }
\end{aligned}
$$

Typical choices for update, aggregate and readout: Multilayer Perceptrons

## MPNN example: GNN 101

* Non-linear activation function $\sigma$ (ReLU, sign, sigmoid, ...)
* $\mathbf{F}_{v o}^{(t)} \in \mathbb{R}^{d}$ denotes embedding of vertex $v$

* Weight matrices $\mathbf{W}_{1}^{(t)} \in \mathbb{R}^{d \times d}$ and $\mathbf{W}_{2}^{(t)} \in \mathbb{R}^{d \times d}$ and bias vector $\mathbf{b} \in \mathbb{R}^{1 \times d}$

Matrix form

$$
\begin{aligned}
& \mathbf{F}_{v_{0}}^{(0)}:=L_{G}(v) \longleftarrow \text { Embedding vertex labels } \\
& \mathbf{F}_{v_{0}}^{(t)}:=\sigma\left(\mathbf{F}_{v 0}^{(t-1)} \mathbf{W}_{1}^{(t)}+\sum_{u \in N_{G}(v)} \mathbf{F}_{u 0}^{(t-1)} \mathbf{W}_{2}^{(t)}+\mathbf{b}^{(t)}\right) \\
& \mathbf{F}^{(t)}:=\sigma\left(\mathbf{F}^{(t-1)} \mathbf{W}_{1}^{(t)}+\mathbf{A} \mathbf{F}^{(t-1)} \mathbf{W}_{2}^{(t)}+\mathbf{B}^{(t)}\right) \quad \begin{array}{c}
\text { Aggregation ov } \\
\text { neighbours }
\end{array}
\end{aligned}
$$

## GNN 101: Graph embedding

* Weight matrix $\mathbf{W} \in \mathbb{R}^{d \times d}$ and and bias vector $\mathbf{b} \in \mathbb{R}^{1 \times d}$

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& \mathbf{F}^{(t)}:=\sigma\left(\sum_{v \in V_{G}} \mathbf{F}^{(L)} \mathbf{W}+\mathbf{b}\right)_{\text {Aggregation over all }} \\
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$$

ERM: Find best parameters $\mathbf{W}_{1}^{(1)}, \ldots, \mathbf{W}_{1}^{(L)}, \mathbf{W}_{2}^{(1)} \ldots, \mathbf{W}_{2}^{(L)}, \mathbf{W}, \mathbf{b}^{(1)}, \ldots, \mathbf{b}^{(L)}, \mathbf{b}$

## Two more examples of MPNNs

* Graph Isomorphism Networks (GIN)

$$
\mathbf{F}_{v 0}^{(t)}:=\operatorname{MLP}^{(t)}\left(\left(1+\epsilon^{(t)}\right) \mathbf{F}_{v 0}^{(t-1)}+\sum_{u \in N_{G}(v)} \mathbf{F}_{u 0}^{(t-1)}\right)
$$

* Graph Convolution Network (GCN)

$$
\mathbf{F}_{v_{0}}^{(t)}:=\operatorname{MLP}^{(t)}\left(\frac{1}{\sqrt{\mid N_{G}(v)+1}} \sum_{u \in N_{G}(v) \cup\{u\}} \frac{1}{\sqrt{\mid N_{G}(u)+1}} \mathbf{F}_{u 0}^{(t-1)}\right)
$$

## MPNNs: Expressive power

What is $\rho$ (MPNNs)?

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Recall: All pairs of graphs $(G, H)$ such that all MPNNs return same graph embedding on both graphs.

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Understanding $\rho$ (MPNNs) translates in understanding power of GNN 101, GCNs, GINs, ....

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Understanding $\rho$ (MPNNs) translates in understanding power of GNN 101, GCNs, GINs, ....

A short detour to graph isomorphism testing

## MPNNs and isomorphic graphs

* Because of invariance: MPNNs embed isomorphic graphs in the same way. That is, if $G \cong H \Rightarrow(G, H) \in \rho($ MPNN $)$
* Can MPNNs embed non-isomorphic graphs differently?

Equivalence class of Isomorphic graphs


## The graph isomorphism problem

Given two graph $G=\left(V_{G}, E_{G}, L_{G}\right)$ and $H=\left(V_{H}, E_{H}, L_{H}\right)$ : are they isomorphic? Or is $\underline{G \cong H}$ ?

* Does there exist a graph isomorphism $\pi: V_{G} \rightarrow V_{H}$ ?
* Theory: computational complexity is open.
* Quasi-polynomial algoritm $n^{\log (n)^{(1)}}$ by László Babai (2016).
* Practice: mostly solvable very fast.


## One-sided test: Colour refinement

Apply heuristic on $G$ and $H$ : If Heuristic say "no" then $G \nsubseteq H$, otherwise we do not know.

* Common heuristic is colour refinement
* In paper 1968 by Boris Weisfeiler and Andrei Leman.



## Colour refinement

* Initial: All vertices have their original colour (label)
* Iteration: Separation of identically coloured vertices based on colour histograms of neighbours.
* Two graphs are non-isomorphic if they have different colour histograms.



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## Color refinement

* Extensively studied in the theoretical computer science community
* Many different characterisations of when two graphs have the same colour histograms (equivalent for colour refinement).
- Successful on random graphs with high probability
* Weak expressive power


## $\rho$ (colour refinement)



* Cannot distinguish d-regular graphs
* Cannot count cycles (triangles)
* Only tree information



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Back to MPNNs

## MPNNs \& Colour refinement

## Theorem (Xu et al. 2019, Morris et al. 2019)

If colour refinement cannot tell two graphs apart then neither can any MPNN!

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## MPNNs \& Colour refinement

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$$
\begin{aligned}
& \text { Color refinement } \\
& \operatorname{cr}^{(0)}(G, v):=\text { Initial label of } v \\
& \operatorname{cr}^{(t)}(G, v):=\operatorname{Hash}\left(\operatorname{cr}^{(t-1)}(G, v),\left\{\left\{\operatorname{cr}^{(t-1)}(G, u) \mid u \in N_{G}(v)\right\}\right\}\right) \\
& \rho(G):=\left\{\left\{\operatorname{cr}(G, v) \mid v \in V_{G}\right\}\right\}
\end{aligned}
$$

$\longrightarrow$ No MPNN can separate these graphs

## MPNNs \& Colour refinement

Recall:


We have just shown: $\rho$ (colour refinement $) \subseteq \rho($ MPNNs $)$
Expressive power of MPNNs is upper bounded by colour refinement

## Lower bound?

* We have seen that MPNNs cannot separate more graphs than colour refinement.
* Can colour refinement separate more graphs than MPNNs?


## Lower bound?

* We have seen that MPNNs cannot separate more graphs than colour refinement.
* Can colour refinement separate more graphs than MPNNs? No!

```
Theorem (Morris ct al. 2019)
There exists a GNN }101\mathrm{ which can embed G and H differently when colour refinement assigns them different colours
```

* The class of MPNNs is as powerful (or weak) as colour refinement


## What else can we say?

$\rho($ colour refinement $)=\rho($ MPNNs $)$

## What else can we say?

$\rho($ colour refinement $)=\rho($ MPNNs $)$


Other - more insightful - characterisations?

## What else can we say?

$$
\rho(\text { colour refinement })=\rho(\mathrm{MPNNs})
$$



Other - more insightful - characterisations?

A detour to homomorphism counts

## Homomorphisms

* Let $P=\left(V_{P}, E_{P}, L_{P}\right)$ and $G=\left(V_{G}, E_{G}, L_{G}\right)$ be graphs.
* A function $h: V_{P} \rightarrow V_{G}$ is a homomorphism if it is edge preserving $(v, w) \in E_{p} \Rightarrow(h(v), h(w)) \in E_{G}$ and label preserving.



## Homomorphism counts

* Define $\underline{\operatorname{HOM}(P, G)}:=\{$ all homomorphisms from $P$ to $G\}$
- Define hom $(P, G):=|\operatorname{HOM}(P, G)|$.




## Homomorphism counts

* Define $\underline{H O M(P, G)}:=\{$ all homomorphisms from $P$ to $G\}$
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\#vertices $=4$


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## Homomorphisms

* Weaker notion than subgraph isomorphism (see later)
* Underlies semantics of many graph query languages
* Algebra of homomorphism counts: A rich and active area of research.


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Back to MPNNs

## MPNNs and hom counts

## Theorem (Dell et al. 2019, Dvorák 2010)

$$
\begin{aligned}
\operatorname{hom}(T, G) & =\operatorname{hom}(T, H) \text { for all trees } T \\
& \text { if and only if }
\end{aligned}
$$

colour refinement cannot distinguish $G$ from $H$.

> Corollary $\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all trees $T$ if and only if no MPNN can distinguish $G$ from $H$.

Follows from $\rho(\mathrm{cr})=\rho(\mathrm{MPNN})$

* MPNNs can only detect tree information from a graph!



## Beyond distinguishing power?

* Logical expressiveness
* Approximation properties (universality)


## Colour refinement (again)

It was mentioned that $\rho$ (colour refinement) has many characterisations.
Of interest is also a logical one, in particular First-order logic with 2 variables and counting quantifiers $\left(C_{2}\right)$.

$$
\varphi(x)=\exists^{\leq 5} y\left(E(x, y) \wedge \exists^{\geq 2} x\left(E(y, x) \wedge L_{a}(x)\right)\right)
$$

Given graph $G$, vertex $v \in V_{G}$ satisfies $\varphi$ :

$$
(G, v) \vDash \varphi
$$

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$$

Given graph $G$, vertex $v \in V_{G}$ satisfies $\varphi$ : It has at most 5 neighbours $(G, v) \vDash \varphi$ each with at least to neighbours labeled "a"

## Colour refinement and $\mathrm{C}_{2}$

## Theorem (Cai et al. 1992)

Two vertices in a graph have the same colour after t iterations of colour refinement if and only if these vertices satisfy the same unary $C_{2}$ formulas of quantifier depth $t$

$$
\rho(\text { colour refinement })=\rho(\mathrm{MPNNs})=\rho\left(\mathrm{C}_{2}\right)
$$

## Which unary $C_{2}$ formulas can MPNNs express?

* Not all: $\varphi(x):=L_{b}(x) \wedge \exists y L_{r}(y)$

I am blue and there exist a red vertex somewhere...

```
H}\mathrm{ can }\mathscr{C}\mathrm{ -express }\Xi\mathrm{ if there exists a }\xi\in\mathscr{H
such that for all }G\in\mathscr{C},\mathbf{v}\in\mp@subsup{V}{G}{p
\xi(G,\mathbf{v})=\Xi(G,\mathbf{v})
```


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\xi(G,\mathbf{v})=\Xi(G,\mathbf{v})
```



Cannot be reached by message passing!

## Which unary $C_{2}$ formulas can MPNNs express?

* Not all: $\varphi(x):=L_{b}(x) \wedge \exists y L_{r}(y)$
* Graded modal logic: syntactical fragment of $C_{2}$ in which quantifiers are of the form $\exists^{\geq N}\left(E(x, y) \wedge \varphi^{\prime}(y)\right)$


## Theorem (Barceló et al. 2020)

Let $\varphi(x)$ be a unary $C_{2}$ formula. Then, $\varphi(x)$ is equivalent to a graded modal logic formula if and only if $\varphi(x)$ is expressible by the class of MPNNs.

$$
\exists \xi \in \text { MPNNs }: \forall G \in \mathscr{G}, \forall v \in V_{G}:(G, v) \vDash \varphi \Leftrightarrow \xi(G, v)=1
$$

## MPNN+: Extended MPNNs

* Can we extend MPNNs such that all $C_{2}$ formulas (including $\left.\varphi(x):=L_{b}(x) \wedge \exists y L_{r}(y)\right)$ can be expressed?
$\xi^{(t)}(G, v):=\operatorname{Upd}^{(t)}\left(\xi^{(t-1)}(G, v), \operatorname{Agg}^{(t)}\left(\left\{\left\{\xi^{(t-1)}(G, v), \xi^{(t)}(G, u) \mid u \in N_{G}(v)\right\}\right\}\right)\right)$
Add global aggregation in every layer
$\xi^{(t)}(G, v):=\operatorname{Upd}^{(t)}\left(\xi^{(t-1)}(G, v), \operatorname{Agg}^{(t)}\left(\left\{\left\{\xi^{(t-1)}(G, v), \xi^{(t)}(G, u) \mid u \in N_{G}(v)\right\}\right\}\right)\right.$

$$
\left.\operatorname{Read}^{(t)}\left(\left\{\left\{\xi^{(t)}(G, u) \mid u \in V_{G}\right\}\right\}\right)\right)
$$

## MPNNs+

## Theorem (Barceló et al. 2020)

Every unary $C_{2}$ formula $\varphi(x)$ is expressible by the class of MPNNs +

* The corresponding colour refinement version is known as the onedimensional Weisfeiler-Leman algorithm or 1-WL

Can MPNN+ express more formulas? Open problem.

$$
\rho(1-\mathrm{WL})=\rho(\mathrm{MPNNs}+)
$$



## Approximation properties

* Equip set of graphs $\mathscr{G}$ with a topology and assume that $\mathscr{H}$ consist of continuous graph embeddings from $\mathscr{G}$ to $\mathbb{R}$.
* Let $\mathscr{C} \subseteq \mathscr{G}$ be a compact set of graphs.

Theorem (Azizian \& Lelarge 2021. G. and Reutter 2022)
If $\mathscr{H}$ is closed under linear combinations and product, then $\mathscr{H}$ can $\mathscr{E}$-approximate any continuous function $\Xi: \mathscr{C} \rightarrow \mathbb{R}$ satisfying

$$
\rho(\mathscr{H}) \subseteq \rho(\{\Xi\}) .
$$

- Can be generalised to general embeddings with output space $\mathbb{R}^{d}$


## MPNNs: Approximation

Theorem (Azizian \& Lelarge 2021, G. and Reutter 2022)
On compact set of graphs, MPNNs can approximate any continuous graph embedding
$\Xi: \mathscr{C} \rightarrow \mathbb{R}$ satisfying $\rho($ colour refinement $) \subseteq \rho(\{\Theta\})$

* We know $\rho($ MPNNs $)=\rho($ colour refinement $)$
* Update functions can be used to approximate product and take linear combinations of MPNNs
* Intricate relation between distinguishing power and approximation properties


## Universality and graph isomorphism

Theorem (Chen et al. (2019)
In order for a class of methods to be able o approximate any (invariant) continuous functions, the class of methods should be able to distinguish any two non-isomorphic graphs.

Proof

- Minimal size $\rho(\mathscr{H}) \subseteq \rho(\{\Xi\})$

$$
(G, H) \in \rho(\mathscr{H}) \Leftrightarrow G \cong H
$$

## Questions?



Feature Augmentation
Boost the expressive power by adding information

## More expressive MPNNs?

## Feature engineering

* Deep learning and MPNNs have replaced "old school" feature engineering approach.

0

* Number of edges
$\longrightarrow$ Number of cycles of length $5 \longrightarrow \mathbb{R}^{d} \longrightarrow$ SVM
* Centrality measures
* MPNNs were supposed to learn such features automatically ...


## Idea \#1: Adding expressive features

Recall:
Theorem
$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all trees $T$ if and only if no MPNN can distinguish $G$ from $H$.

# Idea \#1: Adding expressive features <br> Theorem <br> Recall: $\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all trees $T$ if and only if no MPNN can distinguish $G$ from $H$. 

* What if we add subgraph information before doing messagepassing?

$\uparrow$
More than trees

## Idea \#1: Adding expressive features

## Theorem

Recall: $\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all trees $T$ if and only if no MPNN can distinguish $G$ from $H$.

* What if we add subgraph information before doing messagepassing?


More than trees


## Structural encodings

1.Choose collection of rooted graph patterns/motifs

$$
\mathscr{P}:=\left\{P_{1}^{r}, \ldots, P_{\ell}^{r}\right\}
$$

2.Choose how to match subgraphs in $\mathscr{P}$ with data graph $G$
3.Add count of matches to vertices as extended features.

$$
P^{r}=\left(V_{P}, E_{P},\{r\}\right)
$$


?

## Matches

* Homomorphism: edge preserving
* Subgraph isomorphism: bijection, edge preserving
* Induced subgraph isomorphism: bijection, edge preserving (both ways)

$$
P^{r}=\left(V_{P}, E_{P},\{r\}\right)
$$

$$
\pi: V_{P} \rightarrow V_{S} \subseteq V_{G} \text { containing } v
$$


subiso( $\left.P^{r}, G^{v}\right)$
indsubiso $\left(P^{r}, G^{v}\right)$

## Matches

* Homomorphism: edge preserving
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$$
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$$
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$$



## $\mathscr{P}$-MPNNs

* Add structural encoding as vertex features and run MPNN

$$
\mathscr{P}:=\left\{P_{1}^{r}, \ldots, P_{t}^{r}\right\}
$$

$\mathscr{P}_{-}$MPNNs
$\xi^{(0)}(G, v):=$ Hot-one encoding of label of vertex $v+\operatorname{hom}\left(P_{1}^{r}, G^{v}\right), \ldots, \operatorname{hom}\left(P_{\ell}^{r}, G^{v}\right)$
$\xi^{(t)}(G, v):=\operatorname{Upd}^{(t)}\left(\xi^{(t-1)}(G, v), \operatorname{Agg}^{(t)}\left(\left\{\left\{\xi^{(t-1)}(G, v), \xi^{(t)}(G, \hat{x}), \operatorname{hom}\left(P_{1}^{r}, G^{u}\right), \ldots, \operatorname{hom}\left(P_{\ell}^{r}, G^{u}\right) \mid u \in N_{G}(v)\right\}\right\}\right)\right)$
$\rho(G):=\operatorname{Readout}\left(\left\{\left\{\xi^{(L)}(G, v) \mid v \in V_{G}\right\}\right\}\right)$
hom counts of patterns

* Did we increase expressive power?


## $\mathscr{P}-\mathrm{MPNNs}$

$$
\mathscr{P}=\{0-2\} \begin{aligned}
& \text { (2) } \underbrace{v}_{0}(2)(2) \\
& \underbrace{(2)}_{G_{1}}(2)
\end{aligned}
$$


$H_{1}$

$$
\text { (c) }=\text { hom count }
$$

* We have seen that these graphs equivalent for colour refinement but clearly not for $\mathrm{O}_{0}-\mathrm{MPNNs}$.
* So, increase in power!
* What is their precise expressive power?


## $\mathscr{P}$-MPNNs: Expressive power

## Theorem

$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all $\mathscr{P}$-pattern trees $T$ if and only if no P-MPNN can distinguish $G$ from $H$.

## $\mathscr{P}$-MPNNs: Expressive power

## Theorem

$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all $\mathscr{P}$-pattern trees $T$ if and only if no P-MPNN can distinguish $G$ from $H$.

$$
\mathscr{P}=\left\{\mathcal{O}_{0}\right\}
$$



Take tree: add in each tree vertex copies of rooted patterns

## $\mathscr{P}$-MPNNs: Expressive power

## Theorem

$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all $\mathscr{P}$-pattern trees $T$ if and only if no P-MPNN can distinguish $G$ from $H$.

$$
\mathscr{P}=\left\{\mathcal{O}_{0}\right\}
$$



Take tree: add in each tree vertex copies of rooted patterns

| Set $(\mathcal{F})$ | MAE |
| :--- | :---: |
| None | $0.47 \pm 0.02$ |
| $\left\{C_{3}\right\}$ | $0.45 \pm 0.01$ |
| $\left\{C_{4}\right\}$ | $0.34 \pm 0.02$ |
| $\left\{C_{6}\right\}$ | $0.31 \pm 0.01$ |
| $\left\{C_{5}, C_{6}\right\}$ | $0.28 \pm 0.01$ |
| $\left\{C_{3}, \ldots, C_{6}\right\}$ | $0.23 \pm 0.01$ |
| $\left\{C_{3}, \ldots, C_{10}\right\}$ | $\mathbf{0 . 2 2} \pm \mathbf{0 . 0 1}$ |

## Choice of matching?

* Graph Substructure Networks (GSNs): use subiso counts.



## Choice of matching?

* Graph Substructure Networks (GSNs): use subiso counts.

* Expressive power of GSN? Reduction to homomorphism counts

$$
\operatorname{Spasm}(\because)=\{\because \neg \Delta, \sqcup, \perp \Delta, \backslash\} .
$$



More hom counts needed for same subgraph iso

## GSNs: Expressive power

## Theorem

## If $\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all $\mathscr{P}^{\star}$-pattern trees $T$, then no SGN can distinguish $G$ from $H$.

* A direct characterisation in terms of subiso is also possible.
* The choice of patterns in $\mathscr{P}$ is crucial
* Simple patterns such as cycles and cliques work well.

The larger and complex $\mathscr{P} \Rightarrow$ more complexity counting

$$
\Rightarrow \text { more expressive power }
$$

- $\mathscr{P}^{\star}$-MPNNs
- $\mathscr{P}$-SGNs - $\mathscr{D}_{-M P N N S}$

Complexity


## Idea \#2: (Random) Vertex identifiers

* Message-Passing is only based on vertex features and adjacency information.
* Two different vertices with the same vertex features will be treated the same (if they have the same colour in colour refinement).

What if we add vertex identifiers?

## Vertex identifiers

Self identification: useful for cycle detection


In terms of colour refinement: every vertex has a unique colour

## Logic comes to rescue

Theorem (Cai et al. 1992)
Recall: Two vertices in a graph have the same colour after $t$ iterations of colour refinement if and only if these vertices satisfy the same unary $C_{2}$ formulas of

If every vertex has a unique colour, then can be identified with a $C_{2}$ formula
We can express in $C_{2}$ a formula $\varphi_{G}$ satisfying

$$
H \vDash \varphi_{G} \Longleftrightarrow H \cong G
$$

## Logic comes to rescue

$$
\begin{aligned}
& \alpha_{v}(x): \bigwedge_{c \text { id of } v} \operatorname{Lab}_{c}(x) \wedge \\
& \beta_{v, w}(x, y):= \begin{cases}\alpha_{v}(x) \wedge \alpha_{w}(y) \wedge E(x, y) & \neg \operatorname{Lab}_{c^{\prime}}(x) \\
\alpha_{v}(x) \wedge \alpha_{w}(y) \wedge \neg E(x, y) & (v, w) \notin \in E_{G}\end{cases} \\
& \varphi_{G}:=\bigwedge_{v \in V_{G}}\left(\exists x \alpha_{v}(x) \wedge \neg \exists^{\geq 2} x \alpha_{v}(x)\right) \wedge \bigwedge_{v, w \in V_{G}} \exists x \exists y \beta_{v, w}(x, y)
\end{aligned}
$$

$$
H \vDash \varphi_{G} \Longleftrightarrow H \cong G
$$

## MPNNs+ and vertex ids

## Recall: Theorem

Every $C_{2}$ formula is expressible by the class of MPNNs+

Idea: We use MPNNs+ to express $\varphi_{G}$

$$
\varphi_{G}:=\bigwedge_{v \in V_{G}}\left(\exists x \alpha_{v}(x) \wedge \neg \exists \geq 2 x \alpha_{v}(x)\right) \wedge \bigwedge_{v, w \in V_{G}} \exists x \exists y \beta_{v, w}(x, y)
$$

## MPNNs+ and vertex ids



## MPNNs+ and vertex ids



## MPNNs+ and vertex ids



## MPNNs+ and vertex ids



## MPNNs+ and vertex ids



## rMPNNs+

* How to choose identifiers? Common choice is at random!
* With high probability random features are vertex identifiers


## Theorem

rMPNNs $(+)$ approximate any invariant graph/ vertex embedding with high probability

* Invariance of computed embedding only in expectation!


## Invariance by averaging

* Add vertex identifiers $G \mapsto(G$, id)
* Take embedding method $\chi \in \mathscr{H}$
* All permutation $\pi \in S_{n}$ with $n=\left|V_{G}\right|$
* Average/Aggregate $P=S_{n}$ :

$$
\xi(G):=\frac{1}{|P|} \sum_{\pi \in P} \xi(\pi(G, \mathrm{id}))
$$

$$
\xi(G):=\max _{\pi \in P} \xi(\pi(G, \mathrm{id}))
$$

## Partial averaging, k-rMPNNs+

Loose interpretation of k-CLIP



## Idea \#3: Use global information

* Extract global graph information and use it as positional encodings of vertices
- Spectral information
* Shortest paths (distance information)
* Biconnectivity (connectivity information)


## Spectral graph theory

* Eigenvalues/vector: $\mathbf{M} \cdot \mathbf{v}=\lambda \mathbf{v}$
* For adjacency matrices: Eigenvalues and eigenvectors of Laplacian $\mathbf{L}_{G}=\mathbf{D}_{G}-\mathbf{A}_{G}$
$\left.\begin{array}{rrrr} & L_{G} \\ -1 & -1 & 0 & 0 \\ -1 & -1 & 0 \\ 0 & -1 & 2 & -1\end{array}\right)$
* Laplacian eigenvalues and vectors contain connectivity information * multiplicity 1 st eigenvalue $\sim$ connected components.


## Spectral MPNNs

Add eigenvectors as vertex features


## MPNNs+eig

* Add Laplacian eigenvectors (spectrum) as features.
eig=eigenvalue+eigenvectors

$$
\left.\begin{array}{l}
\text { SpMPNNS } \\
\xi^{(0)}(G, v):=\text { Hot-one encoding of label of vertex } v+\left(\operatorname{eig}_{1}(v), \ldots, \operatorname{eig}_{n}(v)\right) \\
\xi^{(t)}(G, v):=\operatorname{Upd}^{(t)}\left(\xi^{(t-1)}(G, v), \operatorname{Agg}^{(t)}\left(\left\{\left\{\xi^{(t-1)}(G, v), \xi^{(t)}(G, u), \xi^{(0)},\left(\operatorname{eig}_{1}(u), \ldots, \operatorname{eig}_{n}(u)\right) \mid u \in N_{G}(v)\right\}\right\}\right)\right) \\
\rho(G)
\end{array}\right)=\operatorname{Readout}\left(\left\{\left\{\xi^{(L)}(G, v) \mid v \in V_{G}\right\}\right\}\right) .
$$

* Ambiguity in eigenvector selection
* Not permutation invariant.


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\end{array}\right)=\operatorname{Readout}\left(\left\{\left\{\xi^{(L)}(G, v) \mid v \in V_{G}\right\}\right\}\right) .
$$

* Ambiguity in eigenvector selection
* Not permutation invariant.


## Expressive poser of MPNNs+eig

* Are as powerful as MPNNs with revised vertex labels

* Difficult to analyse.


## Spectral invariant

$$
\mathbf{A}=\sum_{\lambda} \lambda \mathbf{P}_{\lambda} \quad \mathbf{P}_{\lambda}=\left(\begin{array}{cccc}
p_{11}^{\lambda} & p_{12}^{\lambda} & \ldots & p_{1 n}^{\lambda} \\
\vdots & \vdots & \ddots & \vdots \\
p_{n 1}^{\lambda} & p_{n 2}^{\lambda} & \ldots & p_{n n}^{\lambda}
\end{array}\right)
$$

Spectral invariant

$$
v \mapsto \operatorname{specinv}(v):=\left(\lambda, p_{v v}^{\lambda},\left\{\left\{p_{v u}^{\lambda} \mid u \in V_{G}\right\}\right\}\right)_{\lambda \in \Lambda}
$$

## Spectral invariant

$$
\left.\begin{array}{rl}
\mathbf{A}= & \underset{\lambda}{\sum_{\lambda} \lambda \mathbf{P}_{\lambda} \quad \mathbf{P}_{\lambda}=\left(\begin{array}{cccc}
p_{11}^{\lambda} & p_{12}^{\lambda} & \cdots & p_{1 n}^{\lambda} \\
\vdots & \vdots & \ddots & \vdots \\
p_{n 1}^{\lambda} & p_{n 2}^{\lambda} & \cdots & p_{n n}^{\lambda}
\end{array}\right)} \underset{v}{\text { Spectral invariant }}
\end{array}\right) \text { Multiset }
$$

## Spectral invariant

$$
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\end{array}\right)} \underset{\stackrel{\text { Spectral invariant }}{v \mapsto \operatorname{specinv}(v)}}{ }:=\left(\lambda, p_{v v}^{\lambda},\left\{\left\{p_{v u}^{\lambda} \mid u \in V_{G}\right\}\right\}\right)_{\lambda \in \Lambda}
\end{array}\right) \text { Multiset }
$$

Graph properties
Number of length 3, 4, or 5 cycles, whether a graph is connected and the number of length k closed walks from any vertex to itself

## Spectral invariant

$$
\left.\begin{array}{rl}
\mathbf{A}= & \sum_{\lambda}^{\sum_{\lambda} \lambda \mathbf{P}_{\lambda} \quad} \quad \mathbf{P}_{\lambda}=\left(\begin{array}{cccc}
p_{11}^{\lambda} & p_{12}^{\lambda} & \cdots & p_{1 n}^{\lambda} \\
\vdots & \vdots & \ddots & \vdots \\
p_{n 1}^{\lambda} & p_{n 2}^{\lambda} & \cdots & p_{n n}^{\lambda}
\end{array}\right) \\
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\end{array}\right) \text { Multiset }
$$

## Graph properties

Number of length 3, 4, or 5 cycles, whether a graph is connected and the number of length k closed walks from any vertex to itself

Beyond 1-WL/Colour Refinement
$\square$

## SpecMPNN

Spectral invariant

$$
v \mapsto \operatorname{specin} v(v):=\left(\lambda, p_{v v}^{\lambda},\left\{\left\{p_{v u}^{\lambda} \mid u \in V_{G}\right\}\right\}\right)_{\lambda \in \Lambda}
$$

Variation used in Signet and BasisNet $\longrightarrow 2$-WL bound

Can be using combination with any MPNN

## Theorem (Seppelt and Rattan (2023)

specMPNN bounded in power by $(1,1)$-WL and strictly lower than 2-WL

## SpecMPNN

Spectral invariant

$$
v \mapsto \operatorname{specin} v(v):=\left(\lambda, p_{v v}^{\lambda},\left\{\left\{p_{v u}^{\lambda} \mid u \in V_{G}\right\}\right\}\right)_{\lambda \in \Lambda}
$$

Variation used in Signet and BasisNet $\longrightarrow$ 2-WL bound

Can be using combination with any MPNN

## Theorem (Seppelt and Rattan (2023)

specMPNN bounded in power by $(1,1)-$ WL and strictly lower than 2-WL

We discuss these WL's later


## Questions?



Turning one graph into many

## General idea

* Colour refinement equivalent graphs may contain colour refinement inequivalent subgraphs.

$$
\begin{gathered}
1-8 \\
1-8
\end{gathered}
$$

* View graphs as a collection of subgraphs then run MPNN


## General idea

* Colour refinement equivalent graphs may contain colour refinement inequivalent subgraphs.
1-5
* View graphs as a collection of subgraphs then run MPNN


## General idea

* Colour refinement equivalent graphs may contain colour refinement inequivalent subgraphs.

* View graphs as a collection of subgraphs then run MPNN


## Subgraph $\mapsto$ Vertex Aggregation



## Vertex $\mapsto$ Subgraph Aggregation



## The subgraph GNN "wave"



## The subgraph GNN "wave"



## Selection policies

DS-GNN - vertex deletion<br>- edge deletion<br>- ego nets<br>- marked ego-nets



ID-GNNs - marked ego-nets
GNNs-AK - ego-nets
k-OSAN - size k subgraph marking

Popular/effective: ego-nets

$$
\begin{aligned}
& i=1-1 \\
& i-1 \\
& i-1
\end{aligned}
$$

## General Subgraph MPNNs

* We discuss an extension of MPNNs called Ordered Subgraph Aggregation Networks
* General enough to capture most existing methods*
* Theoretical results on expressive power of OSANs translate directly to these methods.


## k-OSAN

## Initialisation:

Selection of k tuple of vertices g

$$
\pi(\nu, \underline{\mathbf{g}}):=\mathrm{UPD}_{\pi}(\text { type of } \underline{\mathbf{g}}, v)
$$

Induqed subgraph
Initial labels

$$
\xi^{(0)}(\nu, \mathbf{g}):=\text { UPD }(\text { type of } \mathbf{g}, v)
$$



Only edges adjacent to $v$

$$
D_{V}-D_{v}-D_{v}-D_{v} D_{v}-D_{v}
$$

Label them differently


Learnable function (MLP)

## k-OSAN

Iteration t: run MPNN for each $\mathbf{g}$

$$
\xi^{(t)}(v, \mathbf{g}):=\operatorname{UPD}^{(t)}\left(\xi^{(t)}(v, \mathbf{g}), \operatorname{AGG}^{(t)}\left(\left\{\left\{\xi^{(t)}(u, \mathbf{g}) \mid u \in V_{G} \text { or } N_{G}(v)\right\}\right\}\right)\right.
$$



Subgraph $\mapsto$ vertex Aggregation

$$
\xi(v):=\operatorname{AGG}\left(\left\{\left\{\xi^{(L)}(\nu, \mathbf{g}) \frac{\pi(v, \mathbf{g}) \neq 0\}\})}{\text { Selection policy }}\right.\right.\right.
$$

## $\mathrm{k}^{-\mathrm{OSAN}^{T}}$

Iteration t: run MPNN for each $\mathbf{g}$

$$
\xi^{(t)}(v, \mathbf{g}):=\operatorname{UPD}^{(t)}\left(\xi^{(t)}(v, \mathbf{g}), \operatorname{AGG}^{(t)}\left(\left\{\left\{\xi^{(t)}(u, \mathbf{g}) \mid u \in V_{G} \text { or } N_{G}(v)\right\}\right\}\right)\right.
$$



Vertex $\mapsto$ Subgraph aggregation

$$
\xi(\mathbf{g}):=\operatorname{AGG}\left(\left\{\left\{\xi^{(L)}(v, \mathbf{g}) \mid \pi(v, \mathbf{g}) \neq 0, \nu \in V_{G}\right\}\right\}\right)
$$

## k-OSAN

## Theorem (Oian et al. 2022)

- k-OSANs and k-OSANs ${ }^{t}$ encompass almost all subgraph methods with selection policy involving k vertices.
- Strictly bounded in expressive power by $(k+1)-W L$
- Incomparable to k-WL.
$\mathrm{k}=2$
* if 2-WL cannot distinguish graphs, then neither can 1-OSANs
* 2-WL can distinguish more graphs than 1-OSANs
* There exists graphs than can be distinguished by l-OSANs but not by MPNNs, and vice versa, there exists graphs that can be distinguished by MPNNs but not by l-OSAS


## Subgraph GNNs

* Can always ensure to be strictly more expressive than MPNNs by including original graph in batch.
* Tractability only when easy subgraph policies are used, i.e., leading to a small number (linear) of subgraphs.
* Seems a good balance between complexity and expressiveness



## Characterisation $\rho(k$-OSAN $)$

* To our knowledge no characterisation of the expressive power of subgraph GNNs (and $k$-OSANs in particular) in terms of homomorphism counts is known.
* An exception are the 1-OSANs.


## Characterisation $\rho(1-\mathrm{OSAN})$

* Let $\mathscr{F}$ be the class of all forests (collection of trees)
* Let $\mathscr{F}^{+}$be collection of graphs obtained by
* Taking forest $F \in \mathscr{F}$
* Taking set $\varnothing \neq B \subseteq V_{F}$ of vertices
* Contracting all vertices in $B$ to a single vertex (removing loops and multi edges).


## Characterisation $\rho(1-\mathrm{OSAN})$

* Let $\mathscr{F}$ be the class of all forests (collection of trees)
* Let $\mathscr{F}^{+}$be collection of graphs obtained by
* Taking forest $F \in \mathscr{F}$
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* Contracting all vertices in $B$ to a single vertex (removing loops and multi edges).



## Characterisation $\rho(1-\mathrm{OSAN})$

* Let $\mathscr{F}$ be the class of all forests (collection of trees)
* Let $\mathscr{F}^{+}$be collection of graphs obtained by
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$$
\begin{aligned}
& n-2 \\
& i-1 \\
& i \rightarrow-n
\end{aligned}
$$

## Characterisation $\rho(1-$ OSAN $)$

* Note: $\mathscr{F}^{+}$contains all trees, but also cycles etc.
* Note: treewidth of elements in $\mathscr{F}^{+}$is at most two.

Theorem (Seppelt \& Rattan, 2023)

$$
\begin{aligned}
\operatorname{hom}(F, G)= & \operatorname{hom}(F, H) \text { for all } F \in \mathscr{F}^{+} \\
& \text {if and only if }
\end{aligned}
$$

no l-OSAN can distinguish $G$ from $H$.

* 1-OSANs (and also ID-aware GNNs, ...) have the ability to detect cycles, etc.



## Questions?



# K-dimensional Weisfeiler-Leman 

Boosting expressive power by higher-order message-passing

## Motivation

* We have seen that many graph embedding methods are bounded in expressive power by 1-WL or colour refinement
* To go beyond this, one can manually add more expressive features.
* In the theoretical computer science community, however, higherorder version of 1 -WL have been studied for a long time.
* Why not use these to build more powerful embedding methods?


## More powerful heuristic

Apply heuristic on $G$ and $H$ : If Heuristic say "no" then $G \nsubseteq H$, otherwise we do not know.
$G \cong H$ ?
Colour refinement $\longrightarrow \mathrm{No} \rightarrow G \not \approx H$


## More powerful heuristic

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Colour refinement $\longrightarrow \mathrm{No} \longrightarrow G \not \approx H$ $\stackrel{\downarrow}{1-\mathrm{WL} \longrightarrow \mathrm{No} \longrightarrow G \nsubseteq H}$


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$G \cong H$ ?
Colour refinement $\longrightarrow \mathrm{No} \rightarrow G \nsubseteq H$



## K-dimensional Weisfeiler-Leman

* Initial: Colour k-tuples of vertices according to label, adjacency and equality information.

$$
\Rightarrow \text { Same colour if same induced subgraph }
$$



Neighbours: two k-tuples $\mathbf{v}=\left(v_{1}, \ldots, v_{k}\right)$ and $\mathbf{w}=\left(w_{1}, \ldots, w_{k}\right)$ are i-neighbours if $v_{j}=w_{j}$ for all $j \neq i$

## K-dimensional Weisfeiler-Leman

* Iteration: k-tuple colour depending on colours of i-neighbours.

$$
\begin{aligned}
\mathrm{w}_{k}^{(t+1)}\left(G, v_{1}, \ldots, v_{k}\right):= & \left(\mathrm{w}_{k}^{(t)}\left(G, v_{1}, \ldots, v_{k}\right), M^{(t)}\left(G, v_{1}, \ldots, v_{k}\right)\right) \\
M^{(t)}\left(G, v_{1}, \ldots, v_{k}\right):= & \left(\mathrm{w}_{k+1}^{(0)}\left(v_{1}, \ldots, v_{k}, w\right),\right. \\
& \mathrm{w}_{k}^{(t)}\left(w, v_{2}, \ldots, v_{k}\right), \\
& \vdots \\
& \left.\mathrm{w}_{k}^{(t)}\left(v_{1}, \ldots, v_{k-1}, w\right) \mid w \in V_{G}\right)
\end{aligned}
$$

* Graphs: Histogram of colours $\mathrm{wl}_{k}^{(L)}(G, v, \ldots, v)$ for all $v \in V_{G}$


## Properties of k-WL

## Theorem (Cai et al. (1992)

Theorem (Cai et al. (1992)
Distinguishability of graphs by k-WL corresponds to distinguishability by $(\mathrm{k}+1)$-variable fragment of FO with counting quantifier $\left(C_{k+1}\right)$

* Graphs of size n: Isomorphism problem solved by n-WL
* Large neighbourhoods (nk) and $n^{k}$ tuples


## Characterisations of $\rho(k-W L)$

## Recall: Theorem (Dell et al. 2018, <br> $$
\begin{aligned} & \operatorname{hom}(T, G)=\operatorname{hom}(T, H) \text { for all trees } T \\ & \text { if and only if } \end{aligned}
$$ <br> colour refinement cannot tell apart $G$ from $H$

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colour refinement cannot tell apart $G$ from $H$

Now: Theorem (Dell et al. 2018, ...)
$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all graphs $T$ of tree width $k$ if and only if
$\mathrm{k}-$ WL cannot tell apart $G$ from $H$

## Characterisations of $\rho(k-\mathrm{WL})$

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& \text { if and only if }
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$$

colour refinement cannot tell apart $G$ from $H$

Now:

## Theorem (Dell et al. 2018.

Measures "how far from being a tree"
$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all graphs $T$ of tree width $k$ if and only if
k-WL cannot tell apart $G$ from $H$

## Treewidth

* A k-tree is a graph that can be obtained starting from a $(\mathrm{k}+1)$-clique and then iteratively adding a vertex connected to a k -clique

$$
\mathrm{k}=2
$$

$\square$

## Treewidth

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$$
k=2 \text { i> }
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$$
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\text { I }
\end{gathered}
$$

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$$
\begin{aligned}
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& \text { B シ シ シ }
\end{aligned}
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Treewidth of a graph is smallest k such that the graph is a partial k -tree

## Treewidth

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k=2 リ リ
* A partial k-tree is a subgraph of a k-tree

Treewidth of a graph is smallest k such that the graph is a partial k -tree

* Trees=Treewidth 1


## Treewidth

* Alternative definition in terms of tree decomposition
* If $u$ and $v$ neighbours then there is "bag" containing them both.
* All bags containing a vertex $v$ from a connected subtree.
* Graph has treewidth k if it has a tree decomposition with bags of size $\mathrm{k}+1$.


## Treewidth



## Treewidth



## Treewidth




## Treewidth




## Treewidth




## Question:

- tw(cycle of length $k$ )?
- tw(k-clique)?


## Back to $\mathscr{P}$-MPNNs and $\mathscr{P}$-GSNs

## Back to $\mathscr{P}-M P N N s$ and $\mathscr{P}-G S N s$

## $\mathscr{P}$-MPNNs

$\xi^{(0)}(G, v):=$ Hot-one encoding of label of vertex $v+\operatorname{hom}\left(P^{r}, G^{v}\right), \ldots, \operatorname{hom}\left(P_{\ell}^{r}, G^{v}\right)$
$\xi^{(t)}(G, v):=\operatorname{Upd}^{(t)}\left(\xi^{(t-1)}(G, v), \operatorname{Agg}^{(t)}\left(\left\{\left\{\xi^{(t-1)}(G, v), \xi^{(t)}(G, u), \operatorname{hom}\left(P^{r}, G^{u}\right), \ldots, \operatorname{hom}\left(P_{\ell}^{r}, G^{u}\right) \mid u \in N_{G}(v)\right\}\right\}\right)\right)$ $\rho(G):=\operatorname{Readout}\left(\left\{\left\{\xi^{(L)}(G, v) \mid v \in V_{G}\right\}\right\}\right)$

## Back to $\mathscr{P}-M P N N s$ and $\mathscr{P}-G S N s$

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\rho(G) & :=\operatorname{Readout}\left(\left\{\left\{\xi^{(L)}(G, v) \mid v \in V_{G}\right\}\right\}\right)
\end{aligned}
$$

## Theorem

$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all $\mathscr{P}$-pattern trees $T$, if and only no $\mathscr{P}$-MPNN can distinguish $G$ from $H$.

## Back to $\mathscr{P}$-MPNNs and $\mathscr{P}$-GSNs

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## Theorem

If the patterns $P$ in $\mathscr{P}$ have maximal tree width $k$ then the power of $\mathscr{P}$-MPPNs is bounded by k-WL. Similar result for $\mathscr{P}^{-G S N}$ using $\mathscr{P}^{\star}$-MPNNs.
$\mathscr{P}_{k}$ max tree width k $\mathscr{P}_{k}^{\star}$ max tree width k



## Idea: higher-order GNNs

## Theorem (Dell et al. 2018

$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all graphs $T$ of tree width $k$ if and only if
k-WL cannot tell apart $G$ from $H$

1-WL $\longrightarrow$ MPNNs

## Idea: higher-order GNNs

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$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all graphs $T$ of tree width $k$ if and only if
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$\operatorname{hom}(T, G)=\operatorname{hom}(T, H)$ for all graphs $T$ of tree width $k$ if and only if
k-WL cannot tell apart $G$ from $H$

k-MPNNs will detect more graph
 than MPNNs

## k-Folklore GNNs (k-FGNs)

$$
\xi^{(t)}(G, \underbrace{v_{1}, \ldots, v_{k}}_{k \text {-vertex embedding }}):=\operatorname{MLP}_{1}^{(t)}\left(\sum_{u \in V_{G}} \prod_{i=1}^{k} \operatorname{MLP}_{2}^{(t)}\left(\xi^{(t-1)}\left(G, v_{1}, \ldots, v_{i-1}, u, v_{i+1}, \ldots, v_{k}\right)\right)\right)
$$

Expressive power?

## k-Folklore GNNs (k-FGNs)

$$
\xi^{(t)}\left(G, v_{1}, \ldots, v_{k}\right):=\operatorname{MLP}_{1}^{(t)}\left(\sum_{u \in V_{G}} \prod_{i=1}^{k} \operatorname{MLP}_{2}^{(t)}\left(\xi^{(t-1)}\left(G, v_{1}, \ldots, v_{i-1}, u, v_{i+1}, \ldots, v_{k}\right)\right)\right)
$$

Expressive power?
Theorem (Maron et al. 2019), Azizian and Lelarge 2021)

$$
\rho(k-\mathrm{FGNN})=\rho(k-\mathrm{WL})
$$

## k-GNNs

A simpler architecture:

$$
\begin{gathered}
\xi^{(t)}\left(G, v_{1}, \ldots, v_{k}\right):=\sigma\left(\xi^{(t-1)}\left(G, v_{1}, \ldots, v_{k}\right) \mathbf{W}_{1}^{(t)}+\left(\sum_{i=1}^{k} \sum_{u \in V_{G}} \xi^{(t)}\left(G, v_{1}, \ldots, v_{i-1}, u, v_{i+1}, \ldots, v_{k}\right)\right) \mathbf{W}_{2}^{(t)}\right) \\
\text { Global aggregation }
\end{gathered}
$$

Expressive power?

## k-GNNs

A simpler architecture:

$$
\begin{gathered}
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\end{gathered}
$$

Expressive power?
Theorem (Morris et al. 2019)

$$
\rho(k-\mathrm{GNN})=\rho(k-\mathrm{WL})
$$

## Linear equivariant layers

$$
L: \mathbb{R}^{n^{k}} \rightarrow \mathbb{R}^{n^{\ell}} \text { s.t. } L\left(\mathbf{P}^{t} \mathbf{X P}\right)=\mathbf{P}^{t} L(\mathbf{X}) \mathbf{P} \text { for all permutation matrices } \mathbf{P}
$$

One can find a $\underline{\text { basis } \mathbf{B}_{\gamma}}$ s.t. $\mathbf{L}=\sum a_{\gamma} \mathbf{B}_{\gamma}$

$n=16, k=\ell=2$
$\rightarrow$ Build higher-order GNN using linear equivariant layers

## k-IGNs

## $\xi^{(t)}\left(G, v_{1}, \ldots, v_{k}\right):=\sigma\left(\sum_{\text {Equality types } \sim}^{\left.\sum_{\gamma} \sum_{w_{1}, \ldots, w_{k}} \mathbf{B}_{\gamma} \mathbf{W}_{\gamma}^{(t)} \xi^{(t-1)}\left(G, w_{1}, \ldots, w_{k}\right)+\sum_{\mu} \mathbf{B}_{\mu} \mathbf{W}_{\mu}^{(t)}\right)}\right.$

Theorem (Maron et al. 2019, G. and Reutter 2022)

$$
\rho(k-\mathrm{IGN})=\rho((k-1)-\mathrm{WL})
$$



## Higher-order methods

* Do not scale well, but are expressive
* Do not leverage sparsity of graphs
* Powerful, but leads to overfitting

There are several attempts to make them scalable without sacrificing power.

## "Local" k-GNNs: k-LGNNs



## "Local" k-GNNs: k-LGNNs



## "Local" k-GNNs: k-LGNNs



## k-LGNNs

$$
\xi^{(t)}\left(G, v_{1}, \ldots, v_{k}\right):=\sigma\left(\xi^{(t-1)}\left(G, v_{1}, \ldots, v_{k}\right) \mathbf{W}_{1}^{(t)}+\left(\sum_{i=1}^{k} \sum_{\underline{\left(u, v_{i}\right) \in E_{G}}} \xi^{(t)}\left(G, v_{1}, \ldots, v_{i-1}, u, v_{i+1}, \ldots, v_{k}\right)\right) \mathbf{W}_{2}^{(t)}\right)
$$

Theorem (Morris et al. (2020), G and Reutter (2022)

$$
\rho((k+1)-\mathrm{WL}) \subsetneq \rho(k-\mathrm{LGNN}) \subsetneq \rho(k-\mathrm{WL})
$$

Can detect distance two $(\mathrm{k}+1)$-cliques



# Let' stop filling in the landscape 



## Semi-conclusion

* Expressivity has been an important concept in graph learning since 2019
* Has been pushing forward the area: different techniques to boost power:
* k-WL, feature augmentation, subgraphs, structured modulated message passing, ....
* Expressive models juggle with
* Complexity, overfitting, ...


## Semi-conclusion

* When methods are shown to be powerful: existential proofs.
* No reason that this power is met in practice.
* Also, distinguishing power is necessary but not sufficient in practice...


## Semi-conclusion

* Expressivity has been an important concept in graph learning since 2019
* Has been pushing forward the area: different techniques to boost power:
* k-WL, feature augmentation, subgraphs, structured modulated message passing, ....
* Expressive models juggle with
* Complexity, overfitting, ...


## What to use?

## Subgraph

* Small graphs
- Good
compromise in general


## Feature Augmentation

* Large training datasets
* Invariance not importnat
* Preprocessing ok


## Higher-order

* Graphs are small
* Efficiency not essential
* Expressivity guarantee needed


## Road ahead

Expressiveness

* A lot of recent (2023 progress)
* WL hierarchy needs better reconciliation with practice
* Hom count characterisations
* Relational


## Connection with Learning??

- Optimisation and training unexplored
* Generalisation properties
* Sample efficiency?


#  <br> $x^{2} / 2 p i$ <br> Bounding embedding methods 

An "easy" way to analyse the power of graph embeddings

## How to get k-WL bounds?

Without knowing k-WL?

## Higher-order MPNNs

* They are a generalisation of classical MPNNs.
* They provide a flexible mechanism to describe various graph learning architectures.

Easy way to obtain upper bounds on the expressive power of graph learning architectures.

## Higher-order MPNNs

* Higher-order MPNNs are defined inductively and declaratively.
* We provide syntax and semantics.
* With each higher-order MPNN $\varphi$ we associate:
- A dimension describing the output feature dimension; and
* A set of free variables and we write $\varphi(\mathbf{x})$ with $\mathbf{x}=\left\{x_{1}, \ldots, x_{\ell}\right\}$.


## Higher-order MPNNs

## Higher-order MPNN

Syntax
$\varphi(\mathbf{x})$ of dimension $d$ and free variables $\mathbf{x}=\left\{x_{1}, \ldots, x_{\ell}\right\}$

Higher-order embedding
Semantics

$$
\xi_{\varphi}: \mathscr{G} \rightarrow\left(\mathscr{V}^{\ell} \rightarrow \mathbb{R}^{d}\right):\left(G, v_{1}, \ldots, v_{\ell}\right) \mapsto \mathbb{R}^{d}
$$

## Higher-order MPNNs: Atomic

Atomic higher-order MPNNs: Syntax
Label: $\varphi\left(x_{i}\right):=\operatorname{Lab}_{j}\left(x_{i}\right)$ of dim 1 and free var $x_{i}$
Edge: $\varphi\left(x_{i}, x_{j}\right):=E\left(x_{i}, x_{j}\right)$ of dim 1 , free vars $x_{i}, x_{j}$
Equality: $\varphi\left(x_{i}, x_{j}\right):=\mathbb{1}\left[x_{i}=x_{j}\right]$ of dim 1 , free vars $x_{i}, x_{j}$

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Equality: $\varphi\left(x_{i}, x_{j}\right):=\mathbb{1}\left[x_{i}=x_{j}\right]$ of dim 1, free vars $x_{i}, x_{j}$

## Semantics

$\xi_{\varphi}:\left(v_{1}, v_{2}, \ldots, v_{p}\right) \mapsto j$ th feature of $v_{i}$
$\xi_{\varphi}:\left(v_{1}, v_{2}, \ldots, v_{p}\right) \mapsto \begin{cases}1 & \left(v_{i}, v_{j}\right) \in E \\ 0 & \text { otherwise }\end{cases}$
$\xi_{\varphi}:\left(v_{1}, v_{2}, \ldots, v_{p}\right) \mapsto \begin{cases}1 & v_{i}=v_{j} \\ 0 & \text { otherwise }\end{cases}$

## Higher-order MPNNs: Atomic



## Higher-order MPNNs: Function Application

Function application: Syntax
Let $\varphi_{1}\left(\mathbf{x}_{1}\right), \ldots, \varphi_{\ell}\left(\mathbf{x}_{1}\right)$ be higher-order MPNNs of $\operatorname{dim} d_{1}, \ldots, d_{\ell}$ and free vars $\mathbf{x}_{1}, \ldots, \mathbf{x}_{\ell}$ Let $F: \mathbb{R}^{d_{1}+\cdots+d_{e}} \rightarrow \mathbb{R}^{d}$ be a function. Then,

$$
\varphi(\mathbf{x})=F\left(\varphi_{1}, \ldots, \varphi_{\ell}\right)
$$

is a higher-order MPNN of $\operatorname{dim} d$ and free vars $\mathbf{x}=\mathbf{x}_{1} \cup \cdots \cup \mathbf{x}_{\ell}$

## Higher-order MPNNs: Function Application

## Function application: Syntax

Let $\varphi_{1}\left(\mathbf{x}_{1}\right), \ldots, \varphi_{\ell}\left(\mathbf{x}_{1}\right)$ be higher-order MPNNs of $\operatorname{dim} d_{1}, \ldots, d_{\ell}$ and free vars $\mathbf{x}_{1}, \ldots, \mathbf{x}_{\ell}$ Let $F: \mathbb{R}^{d_{1}+\cdots+d_{e}} \rightarrow \mathbb{R}^{d}$ be a function. Then,


$$
\xi_{\varphi}:\left(v_{1}, \ldots, v_{p}\right) \mapsto F\left(\xi_{\varphi_{1}}\left(v_{1}, \ldots, v_{p}\right), \ldots, \xi_{\varphi_{t}}\left(v_{1}, \ldots, v_{p}\right)\right)
$$

Linear algebra Activation functions Anything you want...

## Higher-order MPNNs: Aggregation

## Aggregation: Syntax

Let $\varphi_{1}\left(\mathbf{x}_{1}, \mathbf{x}_{2}\right)$ and $\varphi_{2}\left(\mathbf{x}_{1}, \mathbf{x}_{2}\right)$ be higher-order MPNNs of $\operatorname{dim} d_{1}$ and $d_{2}$ and free vars $\mathbf{x}_{1}, \mathbf{x}_{2}$. Let $\Theta$ be a function mapping bags of vectors in $\mathbb{R}^{d_{1}}$ to a vector in $\mathbb{R}^{d}$. Then,

$$
\varphi\left(\mathbf{x}_{1}\right)=\operatorname{agg}_{\mathbf{x}_{2}}^{\Theta}\left[\varphi_{1} \mid \varphi_{2}\right]
$$

is a higher-order MPNN of $\operatorname{dim} d$ and free vars $\mathbf{x}_{1}$

## Higher-order MPNNs: Aggregation

## Aggregation: Syntax

Let $\varphi_{1}\left(\mathbf{x}_{1}, \mathbf{x}_{2}\right)$ and $\varphi_{2}\left(\mathbf{x}_{1}, \mathbf{x}_{2}\right)$ be higher-order MPNNs of $\operatorname{dim} d_{1}$ and $d_{2}$ and free vars $\mathbf{x}_{1}, \mathbf{x}_{2}$. Let $\Theta$ be a function mapping bags of vectors in $\mathbb{R}^{d_{1}}$ to a vector in $\mathbb{R}^{d}$. Then,

$$
\varphi\left(\mathbf{x}_{1}\right)=\operatorname{agg}_{\mathbf{x}_{2}}^{\Theta}\left[\varphi_{1} \mid \varphi_{2}\right]
$$

is a higher-order MPNN of $\operatorname{dim} d$ and free vars $x$

## Semantics

$$
\xi_{\varphi}: \mathbf{v} \mapsto \theta\left(\left\{\left\{\xi_{\varphi_{1}}(\mathbf{v}, \mathbf{w}) \mid \xi_{\varphi_{2}}(\mathbf{v}, \mathbf{w}) \neq \mathbf{0}\right\}\right\}\right)
$$

## Higher-order MPNNs: Aggregation

$\Theta$ is e.g., summation and $\varphi_{2}(x, y):=E(x, y)$ and $\varphi_{1}(x, y):=\mathbb{1}[y=y]$ We can count degrees as follows:

$$
\varphi(x)=\operatorname{agg}_{y}^{\text {sum }}[\mathbf{1}[y=y] \mid E(x, y)]
$$

## Expressive Power of k-MPNNs

A higher-order MPNN is called a k-MPNN if it uses at most k variables. $\mathrm{k}-\mathrm{MPNNs}=$ class of $\mathrm{k}-\mathrm{MPNN}$

## Expressive Power of k-MPNNs

A higher-order MPNN is called a k-MPNN if it uses at most k variables. k-MPNNs=class of k-MPNN

## Theorem (G. And Reutter 2022)

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\rho(k-\mathrm{MPNNs})=\rho(k-\mathrm{WL})
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## Expressive Power of k-MPNNs

A higher-order MPNN is called a k-MPNN if it uses at most k variables. k-MPNNs=class of $\mathrm{k}-\mathrm{MPNN}$

## Theorem (G. And Reutter 2022)

$$
\rho(k-\mathrm{MPNNs})=\rho(k-\mathrm{WL})
$$

Take away: Bounding architectures is easy!!
Just write your architecture as higher-order MPNNs
Count variables

## We end with some examples ...

## MPNNs

We define $\varphi^{(0)}\left(x_{1}\right):=\mathbf{1}\left[x_{1}=x_{1}\right]$
Then for $t>0$, we get

$$
\varphi^{(t)}\left(x_{1}\right):=\operatorname{Upd}^{(t)}\left(\varphi^{(t-1)}\left(x_{1}\right), \operatorname{agg}_{x_{2}}^{\Theta^{(t)}}\left[\varphi^{(t-1)}\left(x_{2}\right) \mid E\left(x_{1}, x_{2}\right)\right]\right)
$$

For readout layer, we get

$$
\left.\varphi:=\operatorname{agg}_{x_{1}}^{\Theta}\left[\varphi^{(L)}\left(x_{1}\right) \mid \mathbb{1}\left[x_{1}=x_{1}\right]\right]\right)
$$

2 variables $\mapsto 1$-WL

## Graph Convolutional Networks

Use $D^{-1 / 2}(I+A) D^{-1 / 2}$ as propagation matrix
$\varphi\left(x_{1}\right):=F\left(\operatorname{agg}_{x_{2}}^{\text {sum }}\left[1\left[x_{2}=x_{2}\right] \mid E\left(x_{1}, x\right)\right]\right)$ with $F: \mathbb{R} \rightarrow \mathbb{R}: x \mapsto \frac{1}{\sqrt{1+x}}$
We can use $\psi\left(x_{1}, x_{2}\right):=\times\left(\times\left(\varphi\left(x_{1}\right),+\left(\mathbf{1}\left[x_{1}=x_{2}\right], E\left(x_{1}, x_{2}\right)\right)\right), \varphi\left(x_{2}\right)\right)$ in the MPNN expressions from the previous slide.

2 variables $\mapsto 1$-WL

## Simplified GNNs

* Uses path information $\mathbf{A}^{p} \mathbf{F}^{(0)}$ in a single layer.
*For $p=3$ and for $\varphi^{(0)}\left(x_{1}\right)$ initial feature:
$* \psi\left(x_{1}\right):=\operatorname{agg}_{x_{2}}^{\text {sum }}\left[\operatorname{aggs}_{x_{1}}^{\text {sum }}\left[\operatorname{agg}_{x_{2}}^{\text {sum }}\left[\varphi^{(0)}\left(x_{2}\right) \mid E\left(x_{1}, x_{2}\right)\right] \mid E\left(x_{2}, x_{1}\right)\right] \mid E\left(x_{1}, x_{2}\right)\right]$

2 variables $\mapsto 1$-WL

## Subgraph count GNNs

* Use count of subgraphs to augment MPNNs
* homomorphism count hom $\left(P^{r}, G^{v}\right)$ for rooted motif $P$,
* subgraph iso count sub $\left(P^{r}, G^{\nu}\right)$ for rooted motif $P$
* If motif has tree width k then hom $\left(P^{r}, G^{\nu}\right)$ can be computed using $\mathrm{k}+\mathrm{l}$ variables.
* For example, $(G, v) \mapsto$ hom $\left(G^{v}\right)$ can be expressed as

$$
\begin{array}{r}
\varphi\left(x_{1}\right):=\sum_{x_{2}} \sum_{x_{3}}
\end{array} \begin{array}{r} 
\\
\left(1 x_{1}, x_{2}\right) E\left(x_{1}, x_{3}\right) E\left(x_{2}, x_{3}\right)\left(\mathbb{1}\left[x_{1}=x_{1}\right]-\mathbb{1}\left[x_{1}=x_{2}\right]\right) \\
\left(1\left[x_{1}=x_{1}\right]-\mathbb{1}\left[x_{1}=x_{3}\right]\right)\left(\mathbb{1}\left[x_{1}=x_{1}\right]-\mathbb{1}\left[x_{2}=x_{3}\right]\right)
\end{array}
$$

$$
\mathrm{k}+1 \text { variables } \mapsto \mathrm{k} \text {-WL }
$$

## Subgraph GNNs: vertices

$$
\begin{aligned}
& \text { MPNN MPNN MPNN MPNN MPNN MPNN } \\
& \varphi^{(t)}\left(x_{1}, x_{2}\right):=\operatorname{Upd}^{(t)}\left(\varphi^{(t)}\left(x_{1}, x_{2}\right):=1\left[x_{1}=x_{2}\right]\right. \\
& \left.1, x_{2}\right), \operatorname{agg}_{x_{3}}^{\Theta}\left[\varphi^{(t-1)}\left(x_{1}, x_{3}\right) \mid E\left(x_{2}, x_{3}\right]\right)
\end{aligned}
$$

Cotta et al.: Reconstruction for powerful graph representations (2021)
Bevilacqua et al.: Understanding and extending subgraph GNNS by rethinking their symmetries (2022
Huang et al.: Boosting the cycle counting power of graph neural networks with I2-GNNs (2022)
Papp et al.: DropGNN: Random dropouts increase the expressiveness of graph neural networks. (2021)
Qian et al.: Ordered subgraph aggregation networks. (2022)
You et al.: Identity-aware graph neural networks. (2021)
Zhang and P. Li. Nested graph neural networks (2021)
Zhao et al.: From stars to subgraphs: Uplifting any GNN with local structure awareness (2022)

3 variables $\mapsto 2$-WL

## Subgraph GNNs: edges



$$
\varphi^{(t)}\left(x_{1}, x_{2}, x_{3}\right):=\operatorname{Upd}^{(t)}\left(\varphi^{(t-1)}\left(x_{1}, x_{2}, x_{3}\right), \operatorname{agg}_{x_{4}}^{\Theta}\left[\varphi^{(t-1)}\left(x_{1}, x_{2}, x_{4}\right) \mid E\left(x_{3}, x_{4}\right]\right)\right.
$$

## 4 variables $\mapsto 3-\mathrm{WL}$

## Conclusion

* Takes a bit of practice but easy to get bounds
* Not guaranteed that these bounds are tight: depends on your programming skills in order to reduce number of variables.
* No lower bounds.

Please use it to get bounds!

